Conjecturing, Testing and Reasoning about Programs Workshop on Verification and Program Transformation 27 March 2021

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Introduction: Theory Exploration

- **Discover interesting properties** about programs.
- **Build** type-correct **terms**. (start small, only non-redundant!)
- Test & evaluate on ground values.
- Assemble equations. (send to prover if you want)
- Use as lemmas, rewrite rules, equational specification...



Demo: Prove my homework Or how to automate a CS undergraduate

fun sorted :: "nat list \Rightarrow bool" where "sorted [] = True" "sorted [x] = True" "sorted (x1#x2#xs) = ((x1 \leq x2) \land sorted (x2#xs))"

fun ins :: "nat \Rightarrow nat list \Rightarrow nat list" where "ins x [] = [x]"

fun isort :: "nat list \Rightarrow nat list" where "isort [] = []" "isort (x#xs) = ins x (isort xs)"

theorem my homework: "sorted (isort x)"

- "ins x (y#ys) = (if (x \leq y) then (x#y#ys) else y#(ins x ys))"

Demo: Prove my homework First attempt: call Sledgehammer

sledgehammer

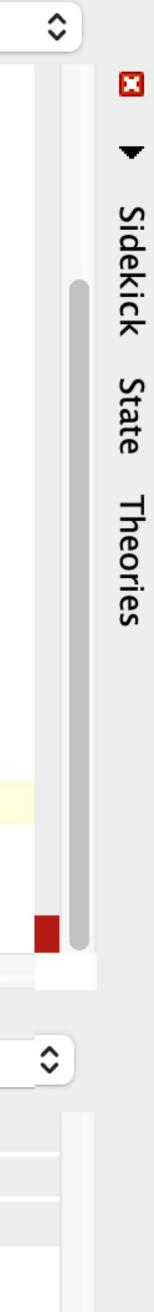
Sledgehammering... "cvc4": Timed out "vampire": Timed out

theorem my_homework: "sorted (isort x)"

Sorted.thy (~/Desktop/)

```
fun sorted :: "nat list \Rightarrow bool"
     where "sorted [] = True"
       "sorted [x] = True"
       "sorted (x1#x2#xs) = ((x1 \leq x2) \land sorted (x2#xs))
   fun ins :: "nat \Rightarrow nat list \Rightarrow nat list"
     where "ins x [] = [x]"
      "ins x (y#ys) = (if (x \leq y) then (x#y#ys) else y#
   fun isort :: "nat list \Rightarrow nat list"
     where "isort [] = []"
      "isort (x#xs) = ins x (isort xs)"
   theorem my_homework: "sorted (isort x)"
     sledgehammer
end
                                                          Upda
                              Proof state
                                         ✓ Auto update
 Sledgehammering...
  "cvc4": Timed out
  "vampire": Timed out
```

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(ins	x ys))"		
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Demo: Prove my homework Six lemmas found and proved

lemma_a: sorted (ins x y) \implies sorted y

 $[lemma_aa: sorted y \implies sorted (ins x y)]$

lemma_ab: ins y (ins x z) = ins x (ins y z)

lemma_ac: sorted (isort x)
(* Needs lemma_aa *)

lemma_ad: "isort (ins x y) = ins x (isort y)"
(* Needs lemma_ab *)

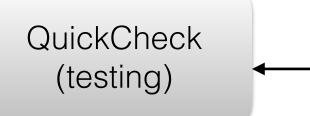
lemma_ae: "isort (isort x) = isort x"
(* Needs lemma_ad *)

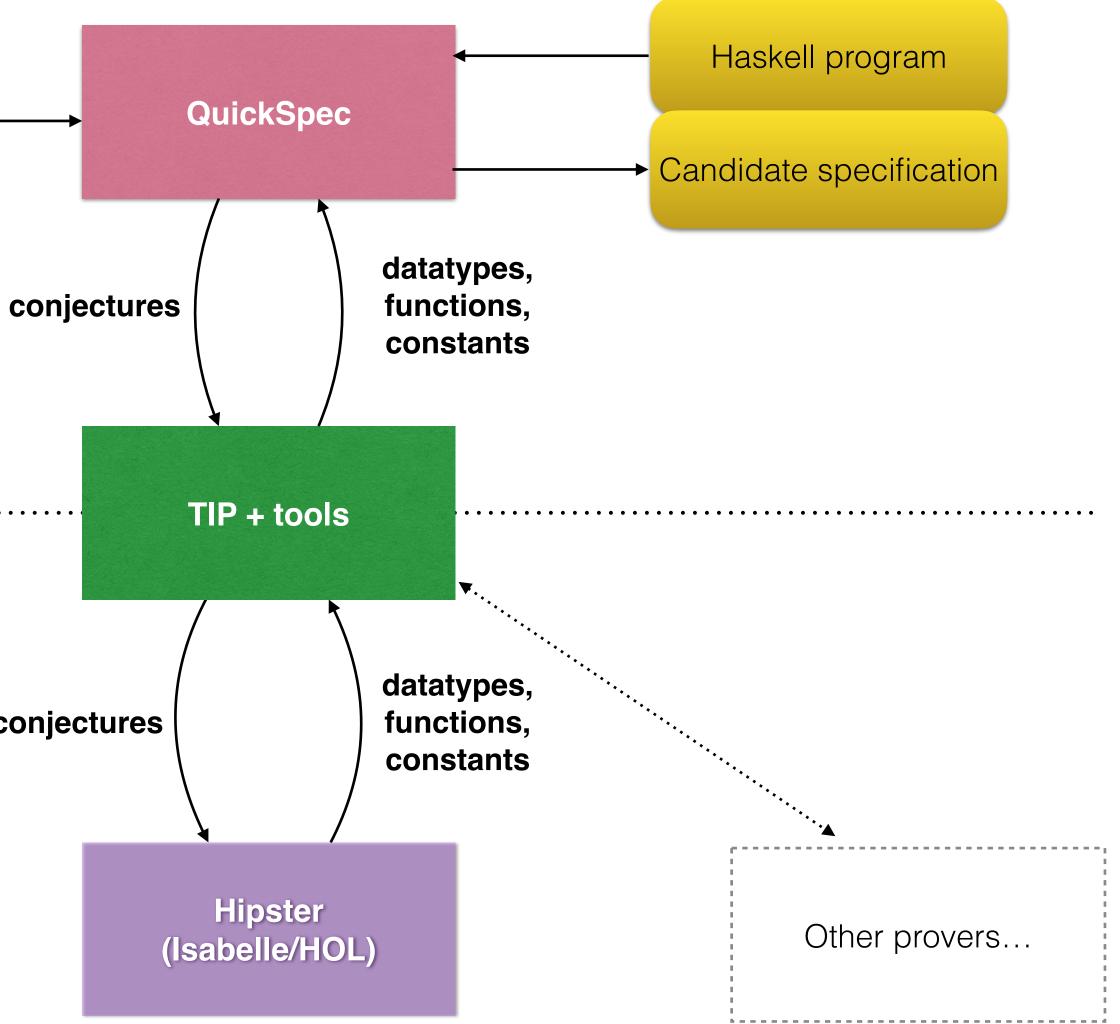
Key lemma

Our homework assignment

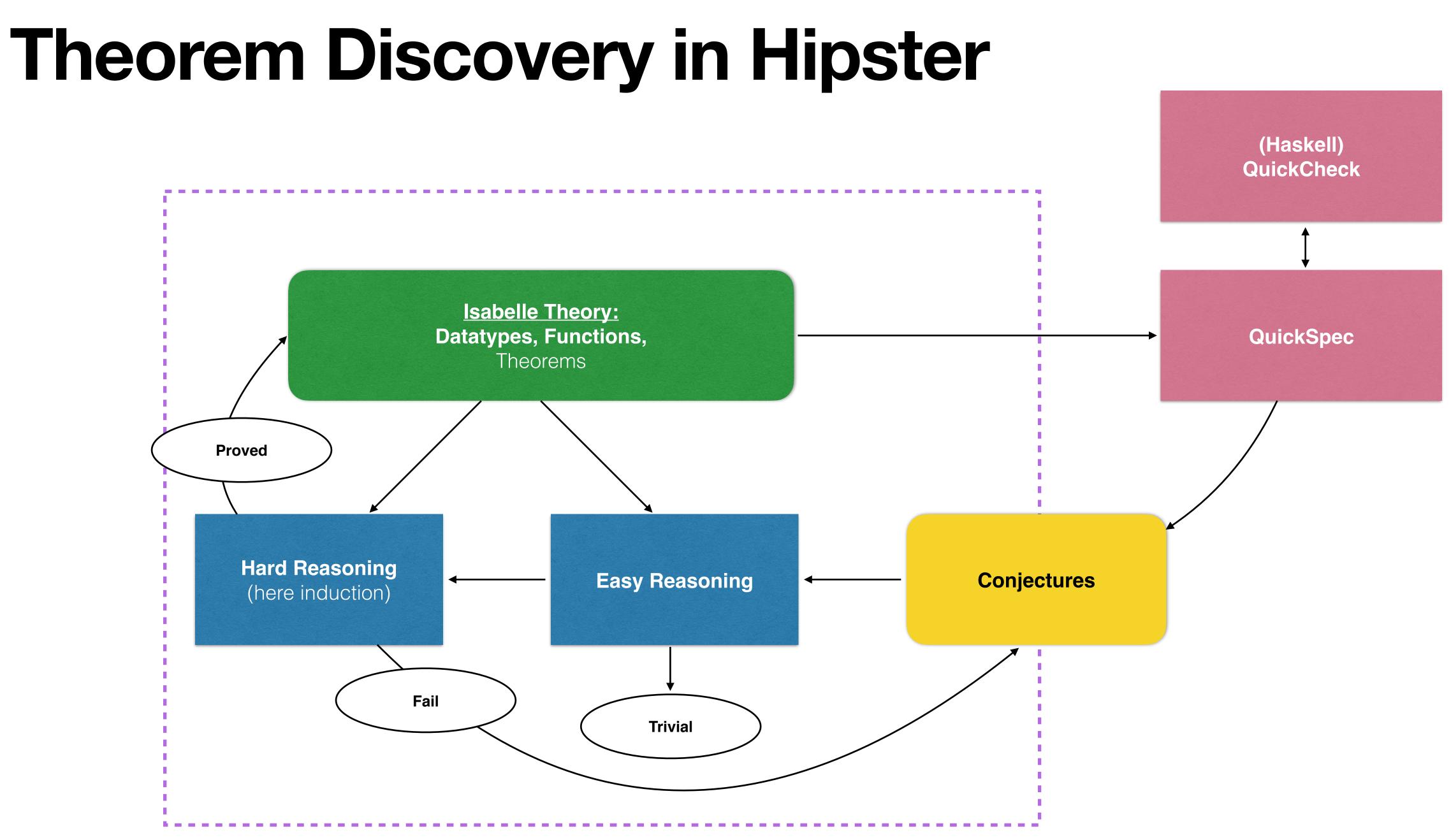
Architecture

.



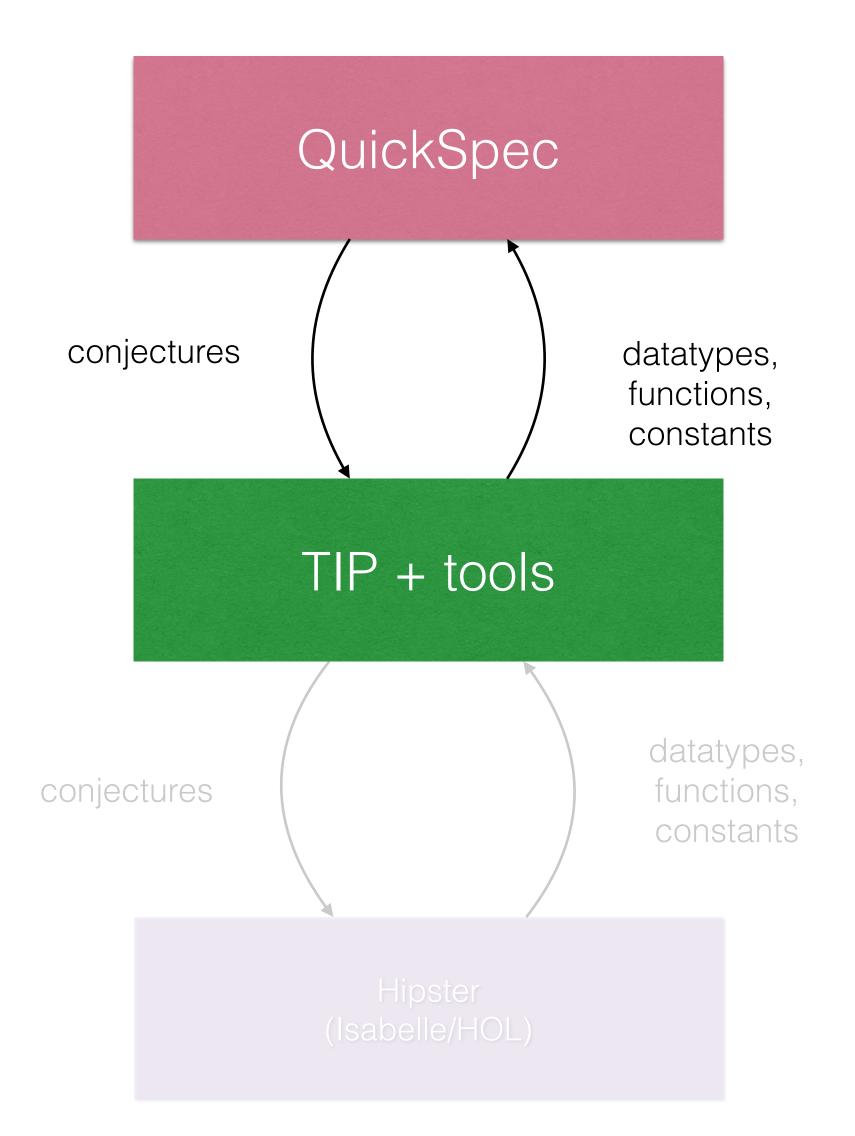


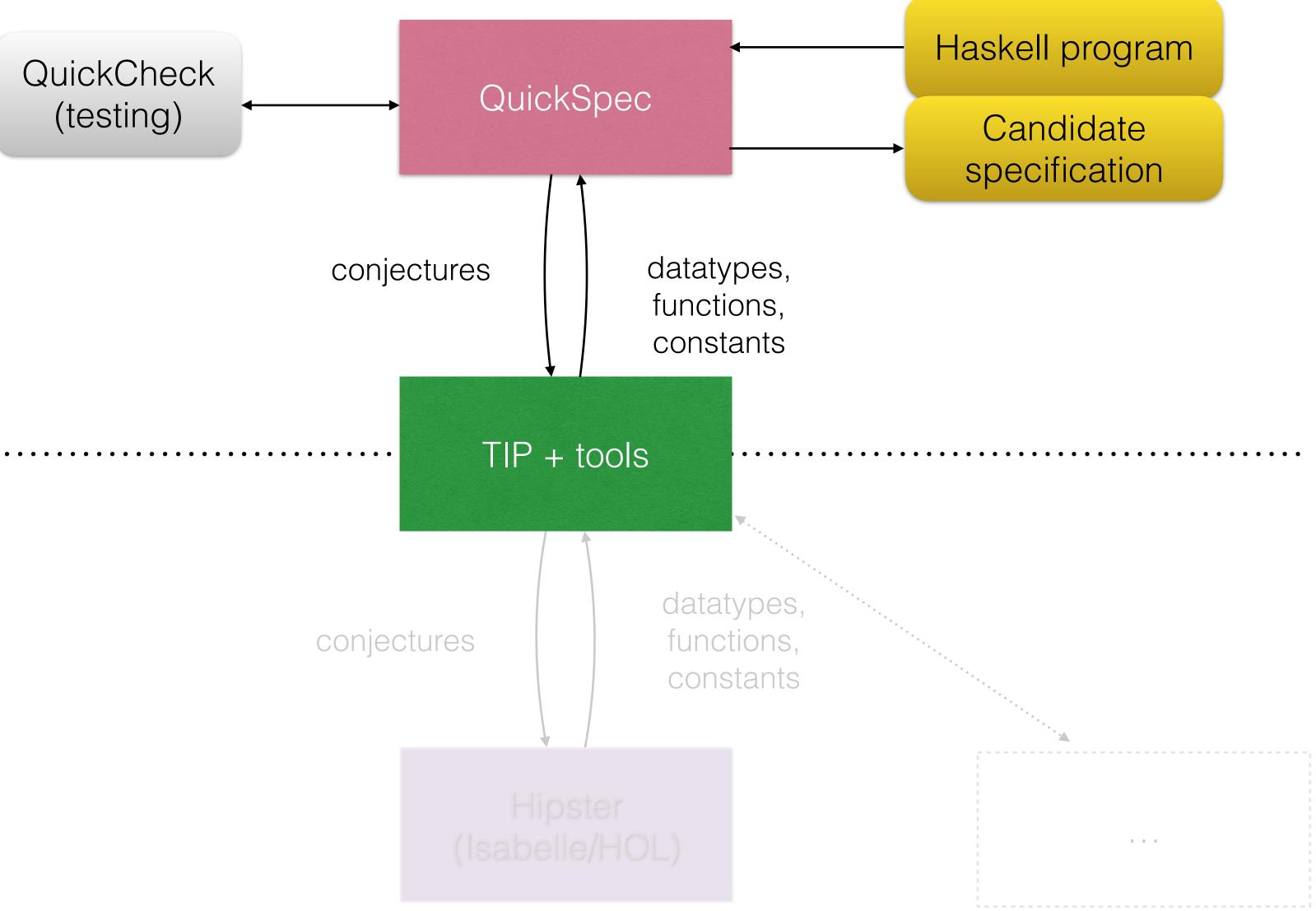
conjectures

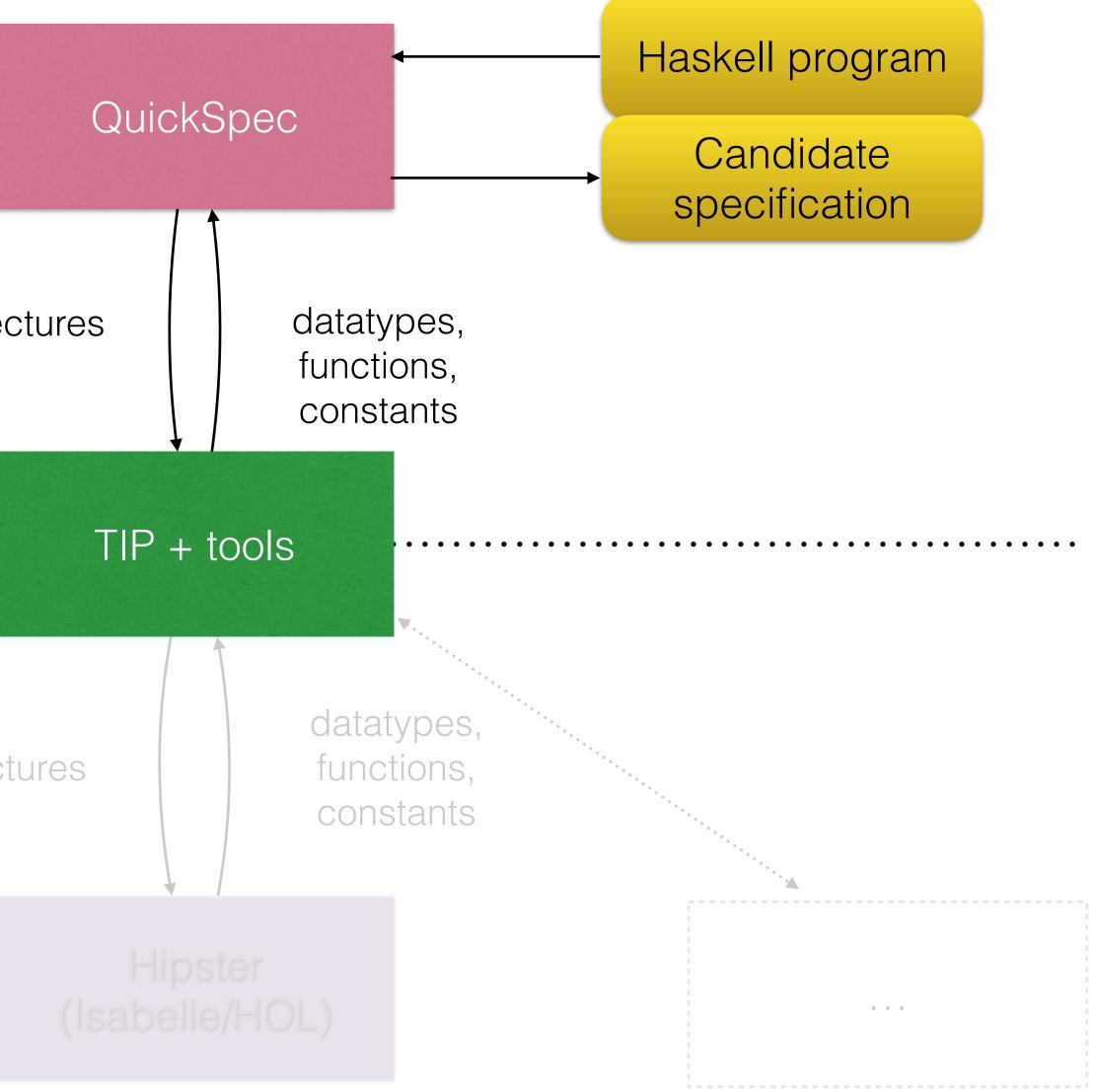


TIP Middle layer

- Intermediate language based on SMT-LIB
 - Datatypes, recursive functions, pattern matching
- Translation tools to various formats, e.g. standard SMT-LIB.
- Libraries for writing prettyprinters and parsers







1. Term generation (small —> large) 2. Testing and evaluation (create equivalence classes) **3. Extract equations. Prune redundant using rewriting.**

ys -> []

XS reverse (reverse xs) sort xs sort (reverse xs) sort (xs++ys)



xs -> [2, 1]

1. Term generation (small —> large) 2. Testing and evaluation (create equivalence classes) 3. Extract equations. Prune redundant using rewriting.

XS reverse (reverse xs) sort xs sort (reverse xs) sort (xs++ys)

[2, 1] reverse (reverse [2, 1]) sort [2, 1] sort (reverse [2, 1]) sort ([2, 1] ++ [])



xs -> [2, 1] ys ->[]

1. Term generation (small —> large) 2. Testing and evaluation (create equivalence classes) 3. Extract equations. Prune redundant using rewriting.





xs -> [2, 1]

	[2, 1]
verse [2, 1])	[2, 1]
	[1, 2]
se [2, 1])	[1, 2]
++ [])	[1, 2]

1. Term generation (small —> large) 2. Testing and evaluation (create equivalence classes) 3. Extract equations. Prune redundant using rewriting.

XS

reverse (reverse xs)

sort xs sort (reverse xs) sort (xs++ys)



xs -> [2, 1, 3] ys -> [3]

1. Term generation (small —> large) 2. Testing and evaluation (create equivalence classes) 3. Extract equations. Prune redundant using rewriting.

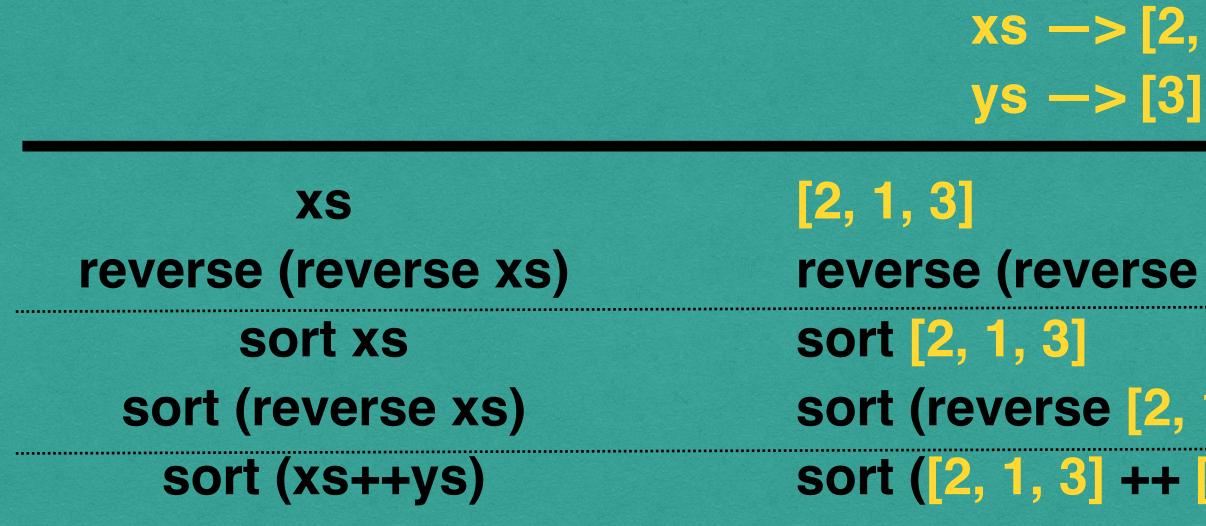




xs -> [2, 1, 3] ys -> [3]

reverse (reverse [2, 1, 3])

1. Term generation (small —> large) 2. Testing and evaluation (create equivalence classes) 3. Extract equations. Prune redundant using rewriting.





xs -> [2, 1, 3]

	[2, 1, 3]
verse [2, 1, 3])	[2, 1, 3]
	[1, 2, 3]
se [2, 1, 3])	[1, 2, 3]
<mark>3] ++ [3])</mark>	[1, 2, 3, 3]

1. Term generation (small —> large) 2. Testing and evaluation (create equivalence classes) 3. Extract equations. Prune redundant using rewriting.

> reverse (reverse xs) = xs sort (reverse xs) = sort xs

reverse (reverse []) = [] reverse (reverse [xs++ys]) = [xs++ys] sort (reverse (sort xs)) = sort (sort xs)

...



jomoa@C20EPLE examples % ./Lists

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Scalability? A stress test

- 30+ functions, large chunk of Haskell's list library.
- Would hit exponential growth of search space...

• But:

- Few equations contain near 30 symbols.
- Interesting properties share similar structure.

```
main = quickSpec [
  con "length" (length :: [A] -> Int),
  con "sort" (sort :: [Int] -> [Int]),
  con "scanr" (scanr :: (A \rightarrow B \rightarrow B) \rightarrow B \rightarrow [A] \rightarrow [B]),
  con "succ" (succ :: Int -> Int),
  con ">>=" ((>>=) :: [A] -> (A -> [B]) -> [B]),
  con "snd" (snd :: (A, B) \rightarrow B),
  con "reverse" (reverse :: [A] -> [A]),
  con "0" (0 :: Int),
  con "," ((,) :: A -> B -> (A, B)),
  con ">=>" ((>=>) :: (A -> [B]) -> (B -> [C]) -> A -> [C]),
  con ":" ((:) :: A -> [A] -> [A]),
  con "break" (break :: (A -> Bool) -> [A] -> ([A], [A])),
  con "filter" (filter :: (A -> Bool) -> [A] -> [A]),
  con "scanl" (scanl :: (B \to A \to B) \to B \to [A] \to [B]),
  con "zipWith" (zipWith :: (A -> B -> C) -> [A] -> [B] -> [C]),
  con "concat" (concat :: [[A]] \rightarrow [A]),
  con "zip" (zip :: [A] -> [B] -> [(A, B)]),
  con "usort" (usort :: [Int] -> [Int]),
  con "sum" (sum :: [Int] -> Int),
  con "++" ((++) :: [A] -> [A] -> [A]),
  con "map" (map :: (A \rightarrow B) \rightarrow [A] \rightarrow [B]),
  con "foldl" (foldl :: (B \to A \to B) \to B \to [A] \to B),
  con "takeWhile" (takeWhile :: (A -> Bool) -> [A] -> [A]),
  con "foldr" (foldr :: (A -> B -> B) -> B -> [A] -> B),
  con "drop" (drop :: Int -> [A] -> [A]),
  con "dropWhile" (dropWhile :: (A -> Bool) -> [A] -> [A]),
  con "span" (span :: (A -> Bool) -> [A] -> ([A], [A])),
  con "unzip" (unzip :: [(A, B)] -> ([A], [B])),
  con "+" ((+) :: Int -> Int -> Int),
  con "[]" ([] :: [A]),
  con "partition" (partition :: (A \rightarrow Bool) \rightarrow [A] \rightarrow ([A], [A])),
  con "fst" (fst :: (A, B) -> A),
  con "take" (take :: Int \rightarrow [A] \rightarrow [A])
```



Current work Scaling exploration

- Limitations: theories with 30+ functions, large terms.
- Use **templates** to skip directly to interesting parts of search space.
- Search only small terms + instances of templates.
- Next: Template mining, applications.

Template-based Theory Exploration: Discovering Properties of Functional Programs by Testing. Sólrún Halla Einarsdóttir, Nicholas Smallbone and Moa Johansson, Proceedings of IFL, to appear 2021.



To conclude...

- Ongoing work get in touch if you want to try it out!
- Hipster/QuickSpec are available for download.

Read more:

- Quick Specifications for the Busy Programmer. Nicholas Smallbone, Moa • Johansson, Koen Claesson and Maximilian Algehed. Journal of Functional Programming, 2017.
- Conference on Interactive Theorem Proving (ITP), p. 1-11, 2017.

Automated Theory Exploration for Interactive Theorem Proving. Moa Johansson.

